

Operating System

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Processes

- Process Concept
- Process Scheduling
- Operations on Processes
- Cooperating Processes
- Interprocess Communication
- Communication in Client-Server Systems

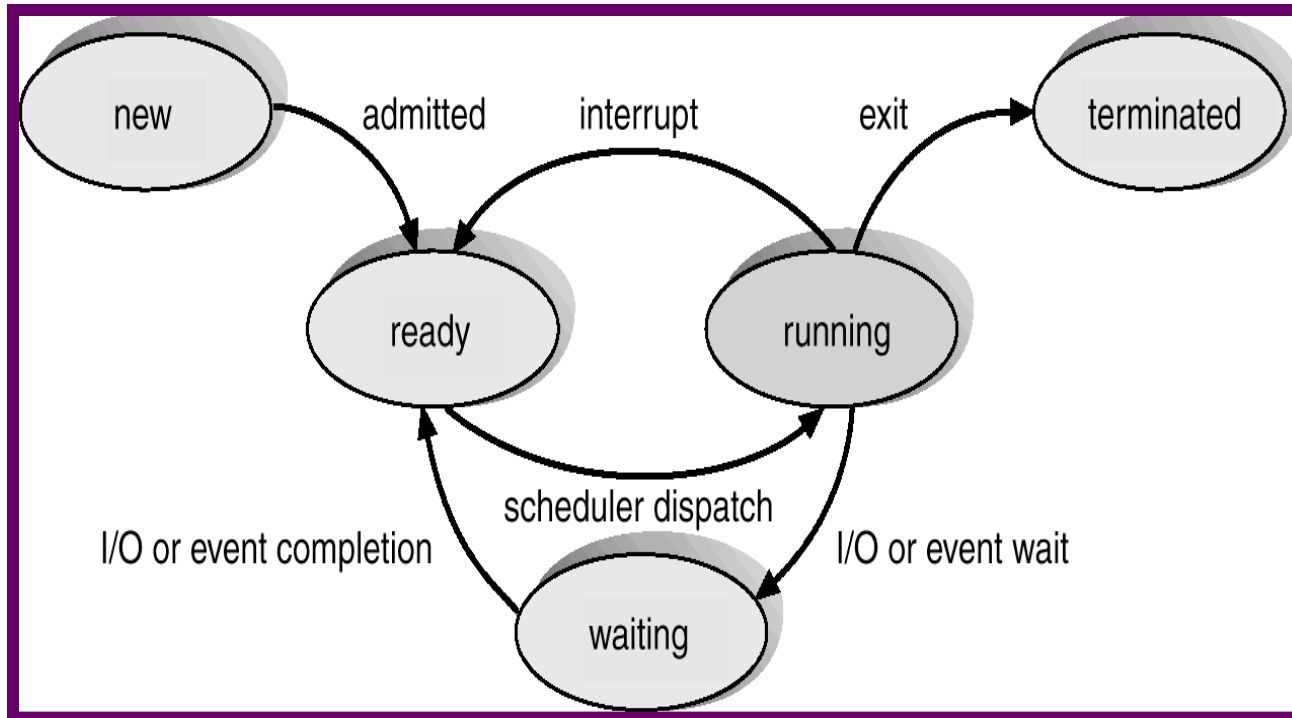
Process Concept

- An operating system executes a variety of programs:
 - } Batch system – jobs
 - } Time-shared systems – user programs or tasks
- Textbook uses the terms *job* and *process* almost interchangeably.
- Process – a program in execution; process execution must progress in sequential fashion.
- A process includes:
 - } program counter
 - } stack
 - } data section

• Process State

- As a process executes, it changes *state*
 - } **new**: The process is being created.
 - } **running**: Instructions are being executed.
 - } **waiting**: The process is waiting for some event to occur.
 - } **ready**: The process is waiting to be assigned to a process.
 - } **terminated**: The process has finished execution.

Diagram of Process State

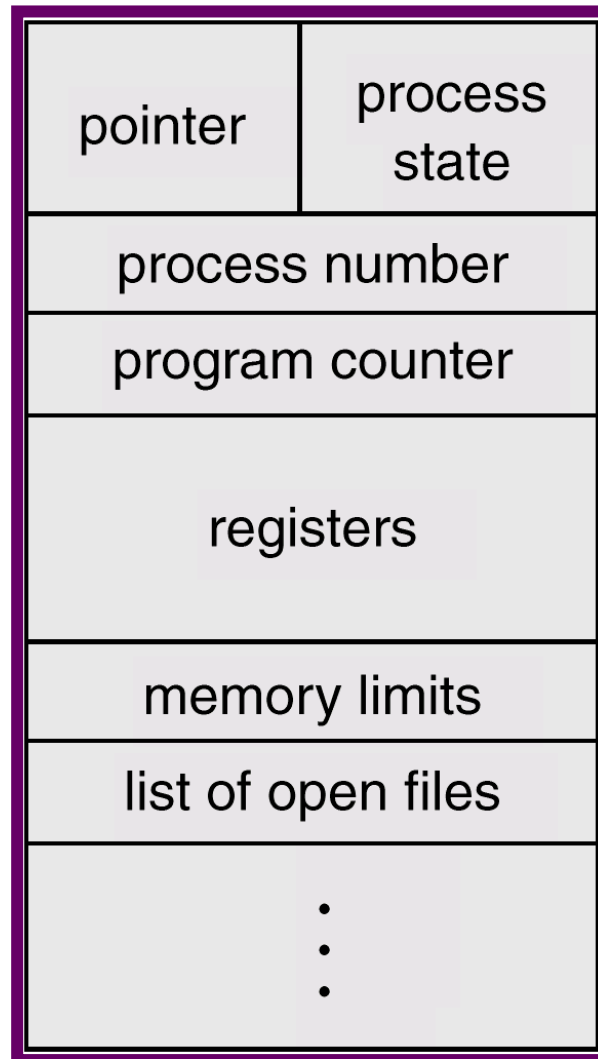


Process Control Block (PCB)

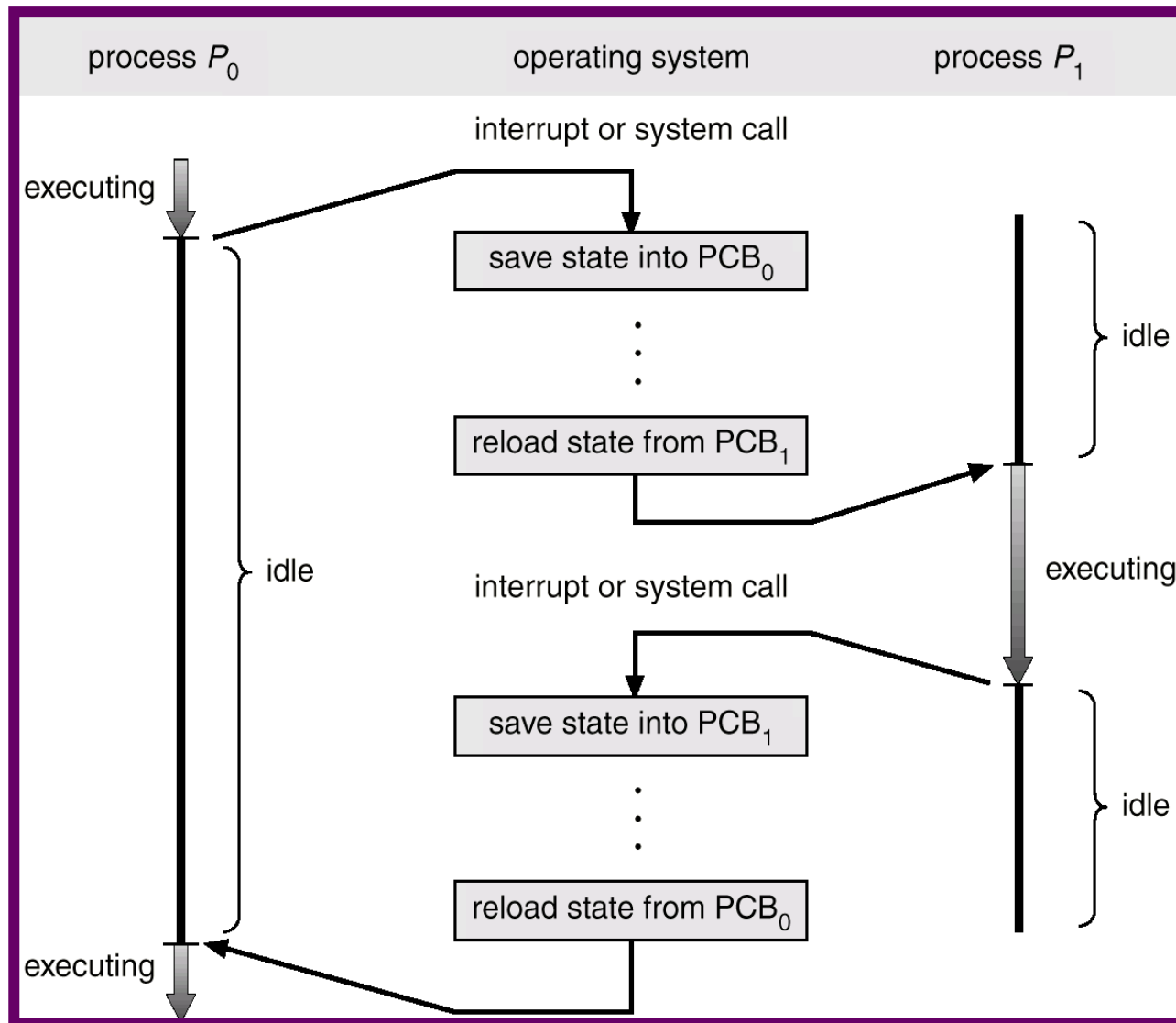
Information associated with each process.

- Process state
- Program counter
- CPU registers
- CPU scheduling information
- Memory-management information
- Accounting information
- I/O status information

Process Control Block (PCB)



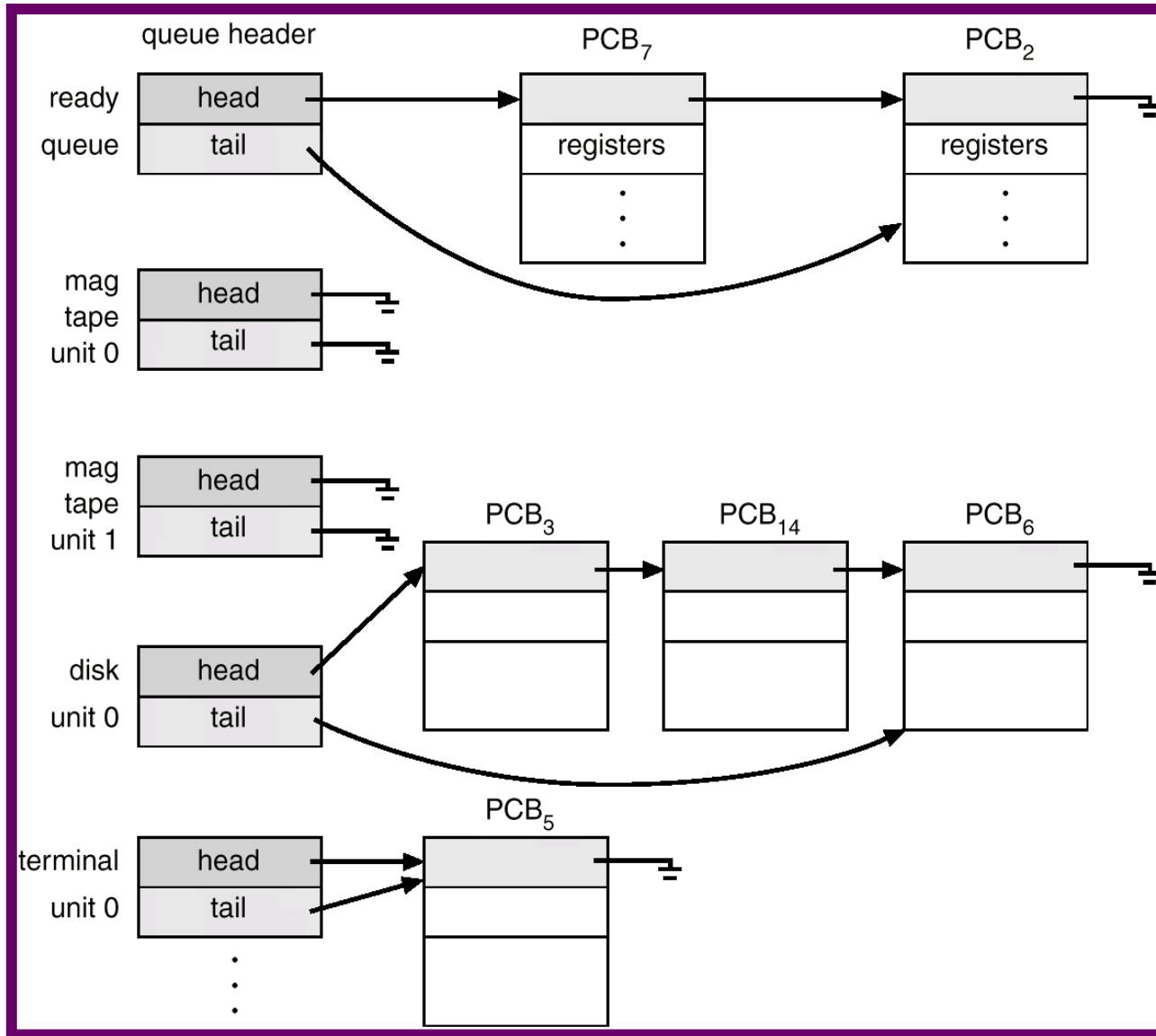
CPU Switch From Process to Process



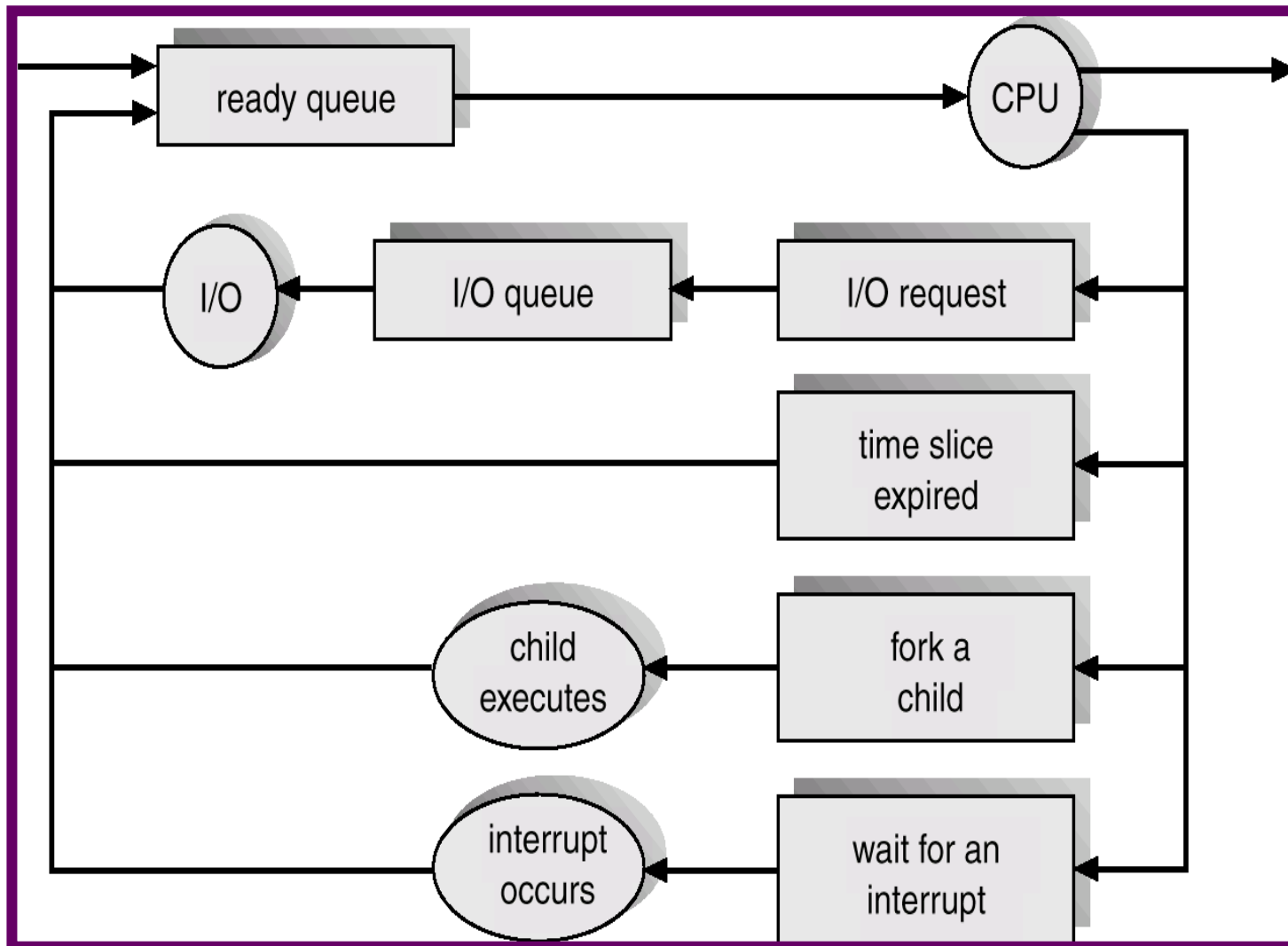
Process Scheduling Queues

- Job queue – set of all processes in the system.
- Ready queue – set of all processes residing in main memory, ready and waiting to execute.
- Device queues – set of processes waiting for an I/O device.
- Process migration between the various queues.

Ready Queue And Various I/O Device Queues



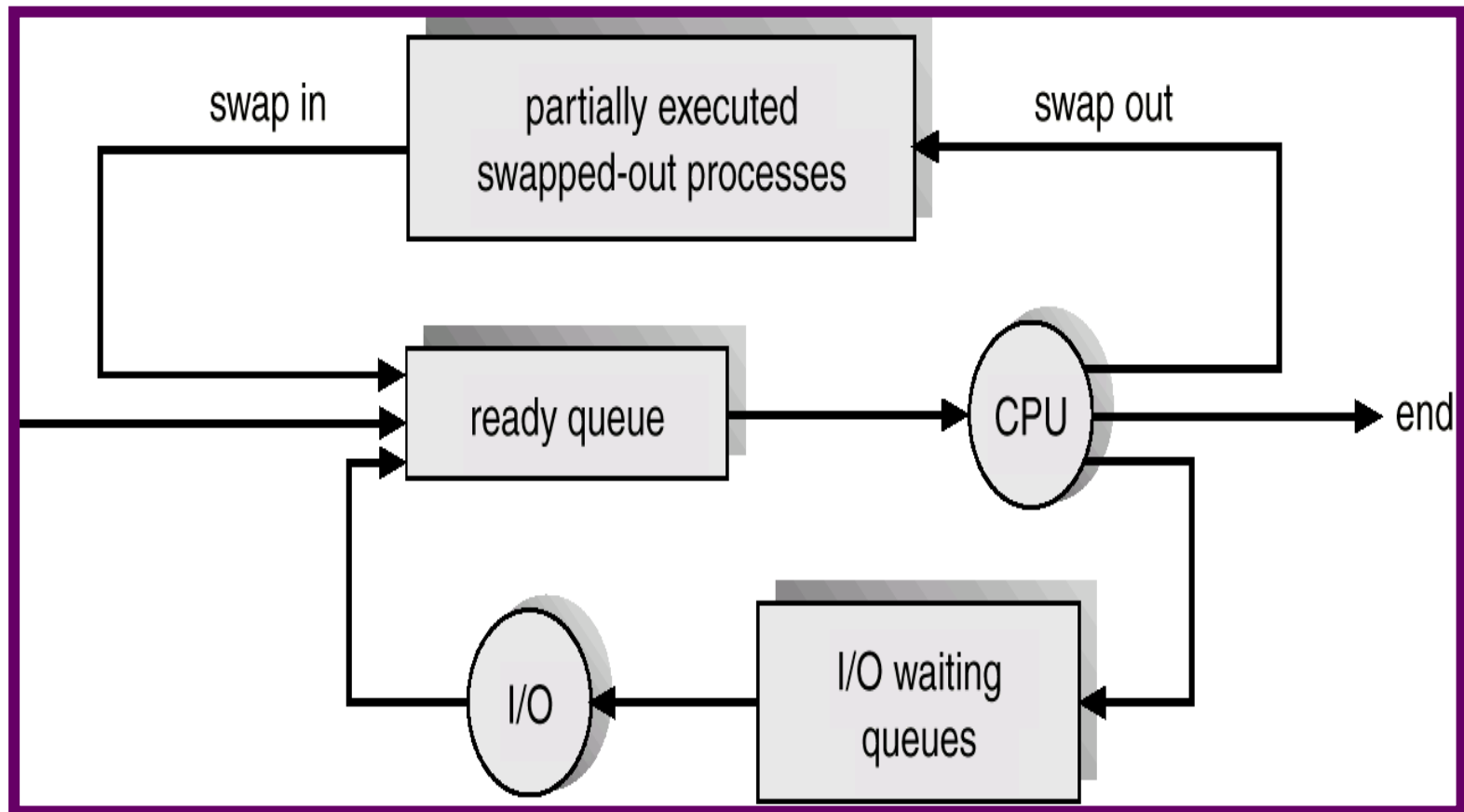
Representation of Process Scheduling



Schedulers

- Long-term scheduler (or job scheduler) – selects which processes should be brought into the ready queue.
- Short-term scheduler (or CPU scheduler) – selects which process should be executed next and allocates CPU.

Addition of Medium Term Scheduling



Schedulers (Cont.)

- Short-term scheduler is invoked very frequently (milliseconds) (must be fast).
- Long-term scheduler is invoked very infrequently (seconds, minutes) (may be slow).
- The long-term scheduler controls the *degree of multiprogramming*.
- Processes can be described as either:
 - *I/O-bound process* – spends more time doing I/O than computations, many short CPU bursts.
 - *CPU-bound process* – spends more time doing computations; few very long CPU bursts.

Context Switch

- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process.
- Context-switch time is overhead; the system does no useful work while switching.
- Time dependent on hardware support.

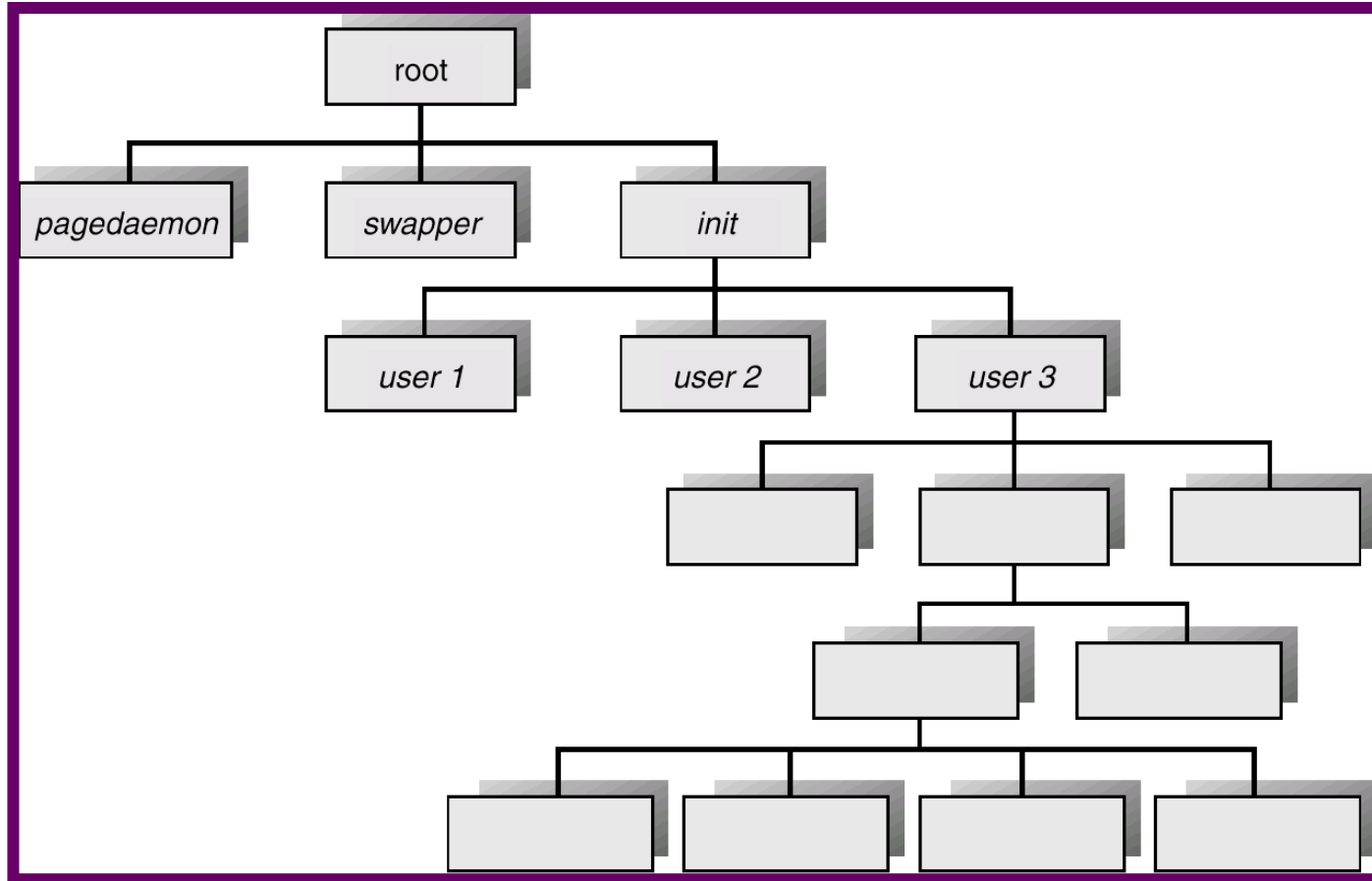
Process Creation

- Parent process create children processes, which, in turn create other processes, forming a tree of processes.
- Resource sharing
 - Parent and children share all resources.
 - Children share subset of parent's resources.
 - Parent and child share no resources.
- Execution
 - Parent and children execute concurrently.
 - Parent waits until children terminate.

• Process Creation (Cont.)

- Address space
 - } Child duplicate of parent.
 - } Child has a program loaded into it.
- UNIX examples
 - } **fork** system call creates new process
 - } **exec** system call used after a **fork** to replace the process' memory space with a new program.

- Processes Tree on a UNIX System



• Process Termination

- Process executes last statement and asks the operating system to decide it (**exit**).
 - › Output data from child to parent (via **wait**).
 - › Process' resources are deallocated by operating system.
- Parent may terminate execution of children processes (**abort**).
 - › Child has exceeded allocated resources.
 - › Task assigned to child is no longer required.
 - › Parent is exiting.
 - Operating system does not allow child to continue if its parent terminates.
 - Cascading termination.

• Cooperating Processes

- *Independent* process cannot affect or be affected by the execution of another process.
- *Cooperating* process can affect or be affected by the execution of another process
- Advantages of process cooperation
 - › Information sharing
 - › Computation speed-up
 - › Modularity
 - › Convenience

• Producer-Consumer Problem

- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process.
 - } *unbounded-buffer* places no practical limit on the size of the buffer.
 - } *bounded-buffer* assumes that there is a fixed buffer size.

- Bounded-Buffer – Shared-Memory Solution

- Shared data

```
#define BUFFER_SIZE 10
typedef struct {
    . . .
} item;
item buffer[BUFFER_SIZE];
int in = 0;
int out = 0;
```

- Solution is correct, but can only use `BUFFER_SIZE-1` elements

- Bounded-Buffer – Producer Process

```
item nextProduced;
```

```
while (1) {  
  while (((in + 1) % BUFFER_SIZE) == out)  
    ; /* do nothing */  
  buffer[in] = nextProduced;  
  in = (in + 1) % BUFFER_SIZE;  
}
```

• Bounded-Buffer – Consumer Process

```
item nextConsumed;
```

```
while (1) {  
    while (in == out)  
        ; /* do nothing */  
    nextConsumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
}
```


• Interprocess Communication (IPC)

- Mechanism for processes to communicate and to synchronize their actions.
- Message system – processes communicate with each other without resorting to shared variables.
- IPC facility provides two operations:
 - } **send**(*message*) – message size fixed or variable
 - } **receive**(*message*)
- If *P* and *Q* wish to communicate, they need to:
 - } establish a *communication link* between them
 - } exchange messages via send/receive
- Implementation of communication link
 - } physical (e.g., shared memory, hardware bus)
 - } logical (e.g., logical properties)

• Indirect Communication

- Operations
 - } create a new mailbox
 - } send and receive messages through mailbox
 - } destroy a mailbox
- Primitives are defined as:
 - send**(*A, message*) – send a message to mailbox *A*
 - receive**(*A, message*) – receive a message from mailbox *A*

• Indirect Communication

- Mailbox sharing
 - } P_1 , P_2 , and P_3 share mailbox A.
 - } P_1 , sends; P_2 and P_3 receive.
 - } Who gets the message?
- Solutions
 - } Allow a link to be associated with at most two processes.
 - } Allow only one process at a time to execute a receive operation.
 - } Allow the system to select arbitrarily the receiver. Sender is notified who the receiver was.

• Synchronization

- Message passing may be either blocking or non-blocking.
- **Blocking** is considered **synchronous**
- **Non-blocking** is considered **asynchronous**
- **send** and **receive** primitives may be either blocking or non-blocking.

• Buffering

- Queue of messages attached to the link; implemented in one of three ways.
 1. Zero capacity – 0 messages
Sender must wait for receiver (rendezvous).
 2. Bounded capacity – finite length of n messages
Sender must wait if link full.
 3. Unbounded capacity – infinite length
Sender never waits.

• Client-Server Communication

- Sockets
- Remote Procedure Calls
- Remote Method Invocation (Java)

• Sockets

- A socket is defined as an *endpoint for communication*.
- Concatenation of IP address and port
- The socket **161.25.19.8:1625** refers to port **1625** on host **161.25.19.8**
- Communication consists between a pair of sockets.

• Socket Communication

